

MUNI

Flatness, weakly-lex colimits and free exact completions

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Exact completion of lex categories

- Exact categories were introduced by Barr as an “ordinary” counterpart of abelian categories.
- Later free exact completions have been introduced:

Definition (Carboni–Magno)

Let \mathcal{C} be lex. The **free (Barr-)exact completion** of \mathcal{C} is an exact category \mathcal{C}_{ex} t.w. $K: \mathcal{C} \hookrightarrow \mathcal{C}_{\text{ex}}$ for which Lan_K induces an equivalence:

$$\text{Lex}(\mathcal{C}, \mathcal{E}) \simeq \text{Ex}(\mathcal{C}_{\text{ex}}, \mathcal{E})$$

for any exact \mathcal{E} .

- \mathcal{C}_{ex} is obtained by freely adding coequalizers of pseudo-equivalence relations to \mathcal{C} .
- $\mathcal{C} \hookrightarrow \mathcal{C}_{\text{ex}} \hookrightarrow [\mathcal{C}^{\text{op}}, \mathbf{Set}]$ closure under finite limits and coequalizers of equivalence relations.

Φ -exact completion of lex categories

Garner–Lack introduce a general notion of Φ -exactness:

$\Phi =$ “class of colimits to which we impose exactness conditions”.

Definition (Garner–Lack)

Let \mathcal{C} be lex. The free Φ -exact completion of \mathcal{C} is a Φ -exact category $\Phi_I \mathcal{C}$ t.w. $K: \mathcal{C} \hookrightarrow \Phi_I \mathcal{C}$ for which Lan_K induces an equivalence:

$$\text{Lex}(\mathcal{C}, \mathcal{E}) \simeq \Phi\text{-Ex}(\Phi_I \mathcal{C}, \mathcal{E})$$

for any Φ -exact \mathcal{E} .

Note: $\mathcal{C} \hookrightarrow \Phi_I \mathcal{C} \hookrightarrow [\mathcal{C}^{\text{op}}, \mathbf{Set}]$ is the closure under finite limits and Φ -lex colimits.

Examples: Regular and Barr-exact categories, (infinitary) lextensive categories, pretopoi, etc.

Problem: does not capture all kinds of free exact completions.

Exact completion of weakly-lex categories

- A diagram $H: \mathcal{D} \rightarrow \mathcal{C}$ has a **weak limit** in \mathcal{C} if there is C t.w. $\delta: \Delta C \rightarrow H$ such that

$$\begin{array}{ccc}
 & \Delta C & \\
 \exists \Delta f \nearrow & & \searrow \delta \\
 \Delta E & \xrightarrow{\quad \forall \eta \quad} & H
 \end{array}$$

- If \mathcal{C} has weak finite limits, then \mathcal{C}_{ex} , obtained by freely adding coequalizers of pseudo-equivalence relations, is exact. (Carboni–Vitale)

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Theorem (Carboni–Vitale)

Let \mathcal{C} be weakly lex and $K: \mathcal{C} \hookrightarrow \mathcal{C}_{ex}$ be the inclusion. Then Lan_K induces an equivalence:

$$\text{Lco}(\mathcal{C}, \mathcal{E}) \simeq \text{Ex}(\mathcal{C}_{ex}, \mathcal{E})$$

for any exact \mathcal{E} .

But what is on the left-hand-side?

Left covering functors

Let $F: \mathcal{C} \rightarrow \mathcal{E}$ be a functor from a weakly lex category \mathcal{C} to a regular category \mathcal{E} .

Definition (Carboni–Vitale/Hu)

We say that F is **left covering** if for any finite diagram $H: \mathcal{D} \rightarrow \mathcal{C}$ and any weak limit $C \in \mathcal{C}$ of H , the comparison map

$$FC \twoheadrightarrow \lim(FH)$$

is a regular epimorphism.

- if \mathcal{C} is lex, then: left covering = lex;
- if $\mathcal{E} = \mathbf{Set}$, then: left covering = flat;

Questions:

- for general \mathcal{C} and \mathcal{E} do we have a “more formal” description?
- can we capture these in the context of Φ -lex colimits?

A notion of flatness

The following are equivalent for $F: \mathcal{C} \rightarrow \mathbf{Set}$:

- ① F is flat (i.e. $\text{El}(F)$ is filtered);
- ② $\text{Lan}_\gamma F: [\mathcal{C}^{\text{op}}, \mathbf{Set}] \rightarrow \mathbf{Set}$ is lex;
- ③ $\text{Lan}_\gamma F: [\mathcal{C}^{\text{op}}, \mathbf{Set}] \rightarrow \mathbf{Set}$ preserves finite limits of representables.

- Replace **Set** with any lex \mathcal{E} ;

Definition

A functor $F: \mathcal{C} \rightarrow \mathcal{E}$, into a lex category \mathcal{E} , is **flat** if and only if for any finite diagram $H: \mathcal{D} \rightarrow \mathcal{C}$, we have

???

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$$\text{Lan}_Y F (\lim YH) \cong \lim (\text{Lan}_Y F \circ YH).$$

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$$\text{colim} \left(\text{El}(\lim YH) \xrightarrow{\pi} \mathcal{C} \xrightarrow{F} \mathbf{Set} \right) \cong \lim FH.$$

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Definition

A functor $F: \mathcal{C} \rightarrow \mathcal{E}$, into a lex category \mathcal{E} , is **flat** if and only if for any finite diagram $H: \mathcal{D} \rightarrow \mathcal{C}$, we have

$$\text{colim} \left(\mathcal{C}/H \xrightarrow{\pi} \mathcal{C} \xrightarrow{F} \mathcal{E} \right) \cong \lim FH.$$

Flatness and free Φ -exact completions

Some properties:

- if $\mathcal{E} = \mathbf{Set}$, then: flat = flat;
- if \mathcal{C} is lex, then: flat = lex;
- if \mathcal{E} is a Grothendieck topos, then: F is flat iff $\text{Lan}_Y F: [\mathcal{C}^{\text{op}}, \mathbf{Set}] \rightarrow \mathcal{E}$ is lex;

Back to Φ -lex colimits. Given a small \mathcal{C} , consider $\Phi_l \mathcal{C}$ to be the closure of \mathcal{C} in $[\mathcal{C}^{\text{op}}, \mathbf{Set}]$ under finite limits and Φ -lex colimits.

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- if $\mathcal{E} = \mathbf{Set}$, then: flat = flat;
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Back to Φ -lex colimits. Given a small \mathcal{C} , consider $\Phi_I \mathcal{C}$ to be the closure of \mathcal{C} in $[\mathcal{C}^{\text{op}}, \mathbf{Set}]$ under finite limits and Φ -lex colimits.

Definition

The inclusion $K: \mathcal{C} \hookrightarrow \Phi_I \mathcal{C}$ exhibits $\Phi_I \mathcal{C}$ as the **free Φ -exact completion of \mathcal{C}** if left Kan extending along K induces an equivalence

$$\text{Flat}(\mathcal{C}, \mathcal{E}) \simeq \Phi\text{-Ex}(\Phi_I \mathcal{C}, \mathcal{E})$$

for any Φ -exact category \mathcal{E} .

The main theorem

Given \mathcal{C} , define

$$\mathcal{C} \subseteq \Phi^\diamond[\mathcal{C}] \subseteq [\mathcal{C}^{\text{op}}, \mathbf{Set}]$$

by adding those M for which M -weakly-lex colimits exist in every Φ -exact \mathcal{E} .

★ In the exact case, objects of $\Phi^\diamond[\mathcal{C}]$ are coequalizers of pseudo equivalence relations between representables.

Theorem

The following are equivalent for a small category \mathcal{C} :

- ① $K: \mathcal{C} \hookrightarrow \Phi_I \mathcal{C}$ exhibits $\Phi_I \mathcal{C}$ as the free Φ -exact completion of \mathcal{C} ;
- ② $\Phi^\diamond[\mathcal{C}] = \Phi_I \mathcal{C}$;
- ③ $\Phi^\diamond[\mathcal{C}]$ has finite limits of diagrams landing in \mathcal{C} .

★ In the exact case, finite limits in $\Phi^\diamond[\mathcal{C}]$ of diagrams landing in \mathcal{C} are weak limits.

Examples

① Φ_{reg} and Φ_{ex} for regular and exact categories;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free regular completion;
- \mathcal{C} has a free exact completion;
- \mathcal{C} is weakly lex.

For such a \mathcal{C} , a functor $F: \mathcal{C} \rightarrow \mathcal{E}$ into a regular category \mathcal{E} is **flat** if and only if it is **left covering**: for any finite diagram $H: \mathcal{D} \rightarrow \mathcal{C}$ and any weak limit $C \in \mathcal{C}$ of H , the comparison map

$$FC \twoheadrightarrow \lim(FH)$$

is a regular epimorphism.

Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{ilext} for infinitary lextensive categories;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free infinitary lextensive completion;
- \mathcal{C} has finite multilimits.

$H: \mathcal{D} \rightarrow \mathcal{C}$ has a **multilimit** in \mathcal{C} if there exists a family of objects $(C_i)_{i \in I}$ in \mathcal{C} together with cones $\delta_i: \Delta C_i \rightarrow H$ for which:

$$\begin{array}{ccc}
 & \Delta C_i & \\
 \exists! i, \exists! \Delta f \nearrow & & \searrow \delta_i \\
 \Delta E & \xrightarrow{\quad \forall \eta \quad} & H
 \end{array}$$

Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{ilext} for infinitary lextensive categories;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free infinitary lextensive completion;
- \mathcal{C} has finite multilimits.

For such a \mathcal{C} , a functor $F: \mathcal{C} \rightarrow \mathcal{E}$ into an infinitary lextensive category \mathcal{E} is **flat** if and only if it is **finitely multicontinuous**: for any finite diagram $H: \mathcal{D} \rightarrow \mathcal{C}$ with multilimit $(C_i)_{i \in I}$ the comparison

$$\sum_{i \in I} FC_i \xrightarrow{\cong} \lim FH$$

is an isomorphism.

Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{illex} for infinitary lexensive categories;
- ③ Φ_{llex} for lexensive categories;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free lexensive completion;
- \mathcal{C} has finite multi-finite limits.

$H: \mathcal{D} \rightarrow \mathcal{C}$ has a **multi-finite limit** in \mathcal{C} if there is a finite family of objects $(C_i)_{i \leq n}$ in \mathcal{C} together with cones $\delta_i: \Delta C_i \rightarrow H$ for which:

$$\begin{array}{ccc}
 & \Delta C_i & \\
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Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{ilext} for infinitary lextensive categories;
- ③ Φ_{lext} for lextensive categories;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free lextensive completion;
- \mathcal{C} has finite multi-finite limits.

For such a \mathcal{C} , a functor $F: \mathcal{C} \rightarrow \mathcal{E}$ into a lextensive category \mathcal{E} is **flat** if and only if for any finite diagram $H: \mathcal{D} \rightarrow \mathcal{C}$ with multi-finite limit $(C_i)_{i \leq n}$ the comparison

$$\sum_{i \leq n} FC_i \xrightarrow{\cong} \lim FH.$$

is an isomorphism.

Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{ilext} for infinitary lextensive categories;
- ③ Φ_{lext} for lextensive categories;
- ④ $\Phi_{pret} = \Phi_{ex} \cup \Phi_{lext}$ for pretopoi;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free pretopos completion;
- \mathcal{C} has finite fc-limits.

$H: \mathcal{D} \rightarrow \mathcal{C}$ has a **fc-limit** in \mathcal{C} if there is a finite family of objects $(C_j)_{j \leq n}$ in \mathcal{C} together with cones $\delta_j: \Delta C_j \rightarrow H$ for which:

$$\begin{array}{ccc}
 & \Delta C_j & \\
 \exists i, \exists \Delta f \nearrow & & \searrow \delta_i \\
 \Delta E & \xrightarrow{\forall \eta} & H
 \end{array}$$

Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{ilext} for infinitary lextensive categories;
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The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free pretopos completion;
- \mathcal{C} has finite fc-limits.

For such a \mathcal{C} , a functor $F: \mathcal{C} \rightarrow \mathcal{E}$ into a pretopos \mathcal{E} is **flat** if and only if for any fc-limit $(C_i)_{i \leq n}$ of a finite diagram H in \mathcal{C} , the comparison

$$\sum_{i \leq n} FC_i \twoheadrightarrow \lim(FH)$$

is a regular epimorphism.

Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{ilext} for infinitary lextensive categories;
- ③ Φ_{lext} for lextensive categories;
- ④ $\Phi_{pret} = \Phi_{ex} \cup \Phi_{lext}$ for pretopoi;
- ⑤ Φ_{gp} of free groupoid actions, for quasi-based categories;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free Φ -exact completion;
- \mathcal{C} has finite polylimits.

polylimits = multilimits but the factorization is unique up to unique automorphism.

$$\begin{array}{ccc} & \Delta \mathcal{C}_i & \\ \exists ! i, \exists ! \cong \Delta f, \nearrow & \searrow \delta_i & \\ \Delta E & \xrightarrow{\forall \eta} & H \end{array}$$

Examples

- 1 Φ_{reg} and Φ_{ex} for regular and exact categories;
- 2 Φ_{ilext} for infinitary lextensive categories;
- 3 Φ_{lext} for lextensive categories;
- 4 $\Phi_{pret} = \Phi_{ex} \cup \Phi_{lext}$ for pretopoi;
- 5 Φ_{gp} of free groupoid actions, for quasi-based categories;

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free Φ -exact completion;
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For such a \mathcal{C} , a functor $F: \mathcal{C} \rightarrow \mathcal{E}$ into a lextensive category \mathcal{E} is **flat** if and only if it is **finitely polycontinuous**.

Examples

- ① Φ_{reg} and Φ_{ex} for regular and exact categories;
- ② Φ_{ilext} for infinitary lexensive categories;
- ③ Φ_{lext} for lexensive categories;
- ④ $\Phi_{pret} = \Phi_{ex} \cup \Phi_{lext}$ for pretopoi;
- ⑤ Φ_{gp} of free groupoid actions, for quasi-based categories;
- ⑥ $\Phi_{\mathbb{D}} = \mathbb{D}$ -filtered diagrams, for a sound class \mathbb{D} .

The following are equivalent for a small Cauchy complete category \mathcal{C} :

- \mathcal{C} has a free Φ -exact completion;
- $\text{Ind}_{\mathbb{D}}(\mathcal{C})$ has finite limits of diagrams in \mathcal{C} .

Thank You