

Bifunctor Thm. and strictification tensor product for double categories with lax double functors

What is a Gray tensor product?

Gray \otimes gives a closed monoidal structure on a category. For 2Cat, the cat. of 2-categories:

$$2\text{Cat}(\mathcal{A} \otimes \mathcal{B}, \mathcal{C}) \cong 2\text{Cat}(\mathcal{B}, \text{Fun}(\mathcal{A}, \mathcal{C})).$$

- ▶ Writing out what a 2-functor $F : \mathcal{B} \rightarrow \text{Fun}(\mathcal{A}, \mathcal{C})$ is,
- ▶ one obtains “quasi-functor of two variables” $H : \mathcal{A} \times \mathcal{B} \rightarrow \mathcal{C}$ defined by relations among $F(B)(A)$, $A \in \mathcal{A}$, $B \in \mathcal{B}$, and
- ▶ concludes which relations should hold in $\mathcal{A} \otimes \mathcal{B}$.

The relation that differs from what holds in $- \times -$ is:

$$(f \otimes B')(A \otimes g) \stackrel{\neq}{\Rightarrow} (A' \otimes g)(f \otimes B).$$

Gray proved that $\mathcal{A} \otimes \mathcal{B}$ yields a monoidal product on 2Cat.

Candidate for Gray \otimes for double cats and lax double functors

- ▶ In [1] the existence of a Gray monoidal structure is proved for strict double categories, and it was fully described by generators and relations in [3].
- ▶ In [4] we defined $\llbracket \mathcal{A}, \mathcal{B} \rrbracket$ to consist of:
 - 0: **lax** double functors
 - 1v: vert. **lax** transf. • 1h: horiz. **oplax** transf.
 - modifications

We characterized a lax d. functor $F : \mathcal{A} \rightarrow \llbracket \mathcal{B}, \mathcal{C} \rrbracket$, got to the notion of **lax double quasi-functor** $H : \mathcal{A} \times \mathcal{B} \rightarrow \mathcal{C}$, and defined $\mathcal{A} \otimes \mathcal{B}$ (in this lax setting).

What is lost in the lax case?

- ▶ $\llbracket -, - \rrbracket$ is NOT a bifunctor
- ▶ associativity constraint is not an isomorphism
- ▶ no enrichment: the composition on $\llbracket \mathcal{A}, \mathcal{B} \rrbracket$ can not be defined (*horizontal* composition and the interchange law of h. o. t. require invertibility of double functors).
- ↪ no Gray monoidal structure.

References

- ▶ G. Böhm: *The Gray Monoidal Product of Double Categories* (2020)
- ▶ P.F. Faul, G. Manuell, J. Siqueira: *2-Dimensional Bifunctor Theorems and Distributive laws* (2021)
- ▶ B. Femić: *Enrichment and internalization in tricategories, the case of tensor categories and alternative notion to intercategories*
- ▶ B. Femić: *Bifunctor Theorem and Gray monoidal structure for double categories with lax double functors*

Lax double quasi-functor

...consists of lax double functors

$$(-, A) : \mathcal{B} \rightarrow \mathcal{C} \quad \text{and} \quad (B, -) : \mathcal{A} \rightarrow \mathcal{C}$$

coinciding on 0-cells, and 2-cells

$$\begin{array}{ccccc} (B, A) \xrightarrow{(g, A)} (B', A) & \xrightarrow{(B', f)} & (B', A') & & (B, A) \xrightarrow{(g, A)} (B', A) \\ \downarrow = & \boxed{(g, f)} & \downarrow = & (u, A) \downarrow & \boxed{(u, f)} \downarrow & (u, A') & (B, U) \downarrow & \boxed{(g, U)} \downarrow & (B', U) & (u, \tilde{A}) \downarrow & \boxed{(u, \tilde{A})} \downarrow & (B, \tilde{A}) \\ (B, A) \xrightarrow{(B, f)} (B, A') & \xrightarrow{(g, A')} & (B', A') & & (B, A) \xrightarrow{(B, f)} (B, A') & \xrightarrow{(g, A')} & (B', A') & & (B, A) \xrightarrow{(B, f)} (B, A') & \xrightarrow{(g, A')} & (B', A') & & (B, A) \xrightarrow{(B, f)} (B, A') & \xrightarrow{(g, A')} & (B', A') \end{array}$$

in \mathcal{C} which satisfy 20 axioms.

Double category isomorphisms

We introduce double cats and construct isos:

$$q\text{-Lax}_{hop}(\mathcal{A} \times \mathcal{B}, \mathcal{C}) \stackrel{*}{\cong} \text{Lax}_{hop}(\mathcal{A}, \llbracket \mathcal{B}, \mathcal{C} \rrbracket)$$

$$q\text{-Lax}_{hop}(\mathcal{A} \times \mathcal{B}, \mathcal{C}) \cong \text{Dbl}_{hop}(\mathcal{A} \otimes \mathcal{B}, \mathcal{C})$$

$$\rightsquigarrow \text{Dbl}_{hop}(\mathcal{A} \otimes \mathcal{B}, \mathcal{C}) \cong \text{Lax}_{hop}(\mathcal{A}, \llbracket \mathcal{B}, \mathcal{C} \rrbracket)$$

- ▶ Hence, there is a natural isomorphism of sets:

$$\text{Dbl}_{st}(\mathcal{A} \otimes \mathcal{B}, \mathcal{C}) \cong \text{Dbl}_{lx}(\mathcal{A}, \llbracket \mathcal{B}, \mathcal{C} \rrbracket).$$

Double categorical Bifunctor Theorem

Passing to **strict** vert. trans. we get a double functor

$$\mathcal{F} : q\text{-Lax}_{hop}^{st}(\mathcal{A} \times \mathcal{B}, \mathcal{C}) \rightarrow \text{Lax}_{hop}(\mathcal{A} \times \mathcal{B}, \mathcal{C}).$$

It restricts to double equivalences:

$$\mathcal{F}' : q\text{-Lax}_{hop}^{st-u}(\mathcal{A} \times \mathcal{B}, \mathcal{C}) \xrightarrow{\cong} \text{Lax}_{hop}^{u-d}(\mathcal{A} \times \mathcal{B}, \mathcal{C})$$

$$\mathcal{F}'' : q\text{-Ps}_{hop}^{st}(\mathcal{A} \times \mathcal{B}, \mathcal{C}) \xrightarrow{\cong} \text{Ps}_{hop}(\mathcal{A} \times \mathcal{B}, \mathcal{C}).$$

“(Un)currying” double functor

Accordingly we get **uncurrying** double functor:

$$\text{Lax}_{hop}(\mathcal{A}, \llbracket \mathcal{B}, \mathcal{C} \rrbracket^{st}) \rightarrow \text{Lax}_{hop}(\mathcal{A} \times \mathcal{B}, \mathcal{C}).$$

It restricts to a double equivalence - **currying** d.f.

$$\text{Lax}_{hop}^{ud}(\mathcal{A} \times \mathcal{B}, \mathcal{C}) \simeq \text{Lax}_{hop}^u(\mathcal{A}, \llbracket \mathcal{B}, \mathcal{C} \rrbracket^{st-u}).$$

Connection to double monads

There are double category isomorphisms:

$$\text{Lax}_{hop}(*, \mathbb{D}) \cong \text{Mnd}(\mathbb{D})$$

$$q\text{-Lax}_{hop}(* \times *, \mathbb{D}) \cong \text{Mnd}(\text{Mnd}(\mathbb{D})).$$

Bifunctor Theorem as a generalization of Beck’s result on the composition of monads:

$$\begin{array}{ccc} q\text{-Lax}_{hop}(* \times *, \mathbb{D}) & \xrightarrow{\mathcal{F}} & \text{Lax}_{hop}(*, \mathbb{D}) \\ \cong \downarrow & & \downarrow \cong \\ \text{Mnd}(\text{Mnd}(\mathbb{D})) & \xrightarrow{\text{Comp}(\mathbb{D})} & \text{Mnd}(\mathbb{D}) \end{array}$$