# Lecture II <br> <br> Introduction to complex networks 

 <br> <br> Introduction to complex networks}

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## Plan of the course

I.

Networks: definitions, characteristics, basic concepts in graph theory
II. Real world networks: basic properties, Models I
III. Models II
IV. Community structure I
V. Community structure II
VI. Dynamic processes in networks

## Two main classes

Natural systems:
Biologic al networks: genes, proteins...
Foodwebs
Social networks

Infra structure networks:
Virtual: web, email, P2P
Physic al: Intemet, power grids, tra nsport...


## Metabolic Networks

## Protein Interactions

Vertices: meta bolites
Edges: chemic al reactions


## Scientific collaboration networks

Vertices: scientists
Edges: co-a uthored pa pers

Weights: depending on

- number of co-a uthored pa pers
- number of authors of each pa per
- number of citations...


## Actors collaboration networks

Vertices: actors
Edges: co-sta rred movies


## World a irport network



## complete IATA database

V = 3100 aiports $E=17182$ weighted edges
$\mathrm{w}_{\mathrm{ij}}$ \#seats/ (time scale)

## Meta-population networks

Each vertex: intemal struc ture
Edges: transport/traffic


## Intemet

- Computers (routers)
- Satellites
- Modems
- Phone cables
- Optic fibers
- EM waves



## Intemet

## Graph representation

different granula rities


Router Level


## Intemet mapping

- continuously evolving a nd growing
- intrinsic heterogeneity
- self-orga nizing
$\longrightarrow$ Largely unknown topology/properties
Mapping projects:
- Multi-probe rec onstruction (router-level): traceroute
- Use of BG P ta bles for the Autonomous System level (doma ins)

Topology and performance mea surements

Netname:
(1717)
as-ebone(3215)
as-telianetse(3301)
bbn/gte(1)
digex(2548)
ebone(3269)
janet(786)
mci(3561)
sprint(1239)
uunet(701)


## The World-Wide-Web

## Virtual network to find and share information <br> - web pages <br> -hyperilinks



## Sa mpling issues

- social networks: va rious samplings/networks
- transportation network: reliable data
- biological networks: inc omplete samplings
- Intemet: va rious (inc omplete) mapping processes
- WWW: regularcrawls
- . .
possibility of introducing bia ses in the measured network characteristics


## Networks cha racteristic s

Networks: of very different origins

## Do they have a nything in common? Possibility to find common properties?

the abstract character of the graph representation a nd graph theory a llow to a nswer....

## Social networks:

## Milg ram's experiment

Milgram, Psych Today 2, 60 (1967)
Dodds et al., Science 301, 827 (2003)


## Small-world properties

## Intemet:



Distribution of distances between two vertices

## Sma ll-world properties

## Average number of vertices within a distance I



## Aiport network

Scientific collaborations

Intemet

## Small-world properties

N vertices, edges with probability
p:
static random graphs

short distances $(\log \mathrm{N})$

## Clustering coeffic ient



Clustering: My friends will know each other with high probability (typic al example: social networks)


# Topological heterogeneity 

Sta tistic a I a na lysis of c entra lity mea sures:
$P(k)=n_{k} / n=$ probability that a rand omly chosen vertex has degree $k$
also: $\mathrm{P}(\mathrm{b}), \mathrm{P}(\mathrm{w}) . .$.

Two broad classes

- homogeneous networks: light ta ils
- heterogeneous networks: skewed, hea vy ta ils


## Airpla ne route network



Netname:
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## Topologic al heterogeneity

Statistic a I a na lysis of centra lity mea sures


## Broad degree distributions

Power-law tails
$P(k) \sim k^{-\gamma}$,
typic ally $2<\gamma<3$

## Topological heterogeneity

 Sta tistic a I a na lysis of centra lity mea sures
linear scale

Poisson
VS.
Power-law
log-scale

## Exp. vs. Scale-Free

## Poisson distribution





Power-law distribution




## Consequences

Power-law tails: $P(k) \sim k^{-\gamma}$
Averag $\langle k\rangle=\int k P(k) d k$
e

$$
\left\langle k^{2}\right\rangle=\int k^{2} P(k) d k \sim k_{c}^{3-\gamma}
$$

$k_{c}=c$ ut-off due to finite-size diverging degree fluctuations for $\gamma<3$

Level of heterogeneity:


## Other heterogeneity levels



## Other heterogeneity levels



## Betweenness c entra lity

## Clustering a nd correlations


non-trivial structures

## Real networks: summa ry!

|  | Network | Type | 7 | $m$ | $z$ | $\ell$ | $\alpha$ | $C^{(1)}$ | $C^{(2)}$ | r | Ref(s). |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | film actors compary directors math coant horship physics coan thorship biology coauthorship telephone call graph email measages email address books student relationships sexual contacts | undirected | 449913 | 25516482 | 113.43 | 3.48 | 2.3 | 0.20 | 0.78 | 0.208 | [20, 415] |
|  |  | undirected | 7673 | 55392 | 14.44 | 4.60 | - | 0.59 | 0.88 | 0.276 | [105, 322] |
|  |  | undirected | 253339 | 496489 | 3.92 | 7.57 | - | 0.15 | 0.34 | 0.120 | [107, 181] |
|  |  | undirected | 52909 | 245300 | 9.27 | 6.19 | - | 0.45 | 0.56 | 0.363 | $[310,312]$ |
|  |  | undirected | 1520251 | 11803064 | 15.53 | 4.92 | - | 0.088 | 0.60 | 0.127 | [310, 312] |
|  |  | undirected | 47000000 | 80000000 | 3.16 |  | 2.1 |  |  |  | [8,9] |
|  |  | directed | 59912 | 86300 | 1.44 | 4.95 | 1.5/2.0 |  | 0.16 |  | [136] |
|  |  | directed | 16881 | 57029 | 3.38 | 5.22 | - | 0.17 | 0.13 | 0.092 | [320] |
|  |  | undirected | 573 | 477 | 1.66 | 16.01 | - | 0.005 | 0.001 | -0.029 | [45] |
|  |  | undirected | 2810 |  |  |  | 3.2 |  |  |  | [264, 265] |
|  | WWW nd. edu | directed | 269504 | 1497135 | 5.55 | 11.27 | 2.1/2.4 | 0.11 | 0.29 | -0.067 | [14, 34] |
|  | WWW Altavista | directed | 203549046 | 2130000000 | 10.46 | 16.18 | 2.1/2.7 |  |  |  | [74] |
|  | citation network | directed | 783339 | 6716198 | 8.57 |  | 3.0/- |  |  |  | [350] |
|  | Roget's Thesaurus | directed | 1022 | 5103 | 4.99 | 4.87 | - | 0.13 | 0.15 | 0.157 | [243] |
|  | word co-occurrence | undirected | 460902 | 17000000 | 70.13 |  | 2.7 |  | 0.44 |  | [119, 157] |
|  | Internet | undirected | 10697 | 31992 | 5.98 | 3.31 | 2.5 | 0.035 | 0.39 | -0.189 | [86, 148] |
|  | power grid | undirected | 4911 | 6594 | 2.67 | 18.99 | - | 0.10 | 0.080 | -0.003 | [415] |
|  | train routes | undirected | 587 | 19603 | 66.79 | 2.16 | 1.0/1. |  | 0.69 | -0.033 | [365] |
|  | software puckages | directed | 1439 | 1723 | 1.20 | 2.42 | 1.6/1.4 | 0.070 | 0.082 | -0.016 | [317] |
|  | software classes | directed | 1377 | 2213 | 1.61 | 1.51 | - | 0.033 | 0.012 | -0.119 | [394] |
|  | electronic circuits | undirected | 24097 | 53248 | 4.34 | 11.05 | 3.0 | 0.010 | 0.030 | -0.154 | [155] |
|  | peer-to-peer network | undirected | 880 | 1296 | 1.47 | 4.28 | 2.1 | 0.012 | 0.011 | -0.366 | [6, 353] |
|  | metabolic network | undirected | 765 | 3686 | 9.64 | 2.56 | 2.2 | 0.090 | 0.67 | -0.240 | [213] |
|  | protein interactions | undirected | 2115 | 2240 | 2.12 | 6.80 | 2.4 | 0.072 | 0.071 | -0.156 | [211] |
|  | marine food web | directed | 135 | 598 | 4.43 | 2.05 | - | 0.16 | 0.23 | -0.263 | [203] |
|  | freshwater food web | directed | 92 | 997 | 10.84 | 1.90 | - | 0.40 | 0.48 | -0.326 | [271] |
|  | neural network | directed | 307 | 2359 | 7.68 | 3.97 | - | 0.18 | 0.28 | -0.226 | [415, 420] |

## C omplex networks

## Complex is not just "c omplic ated"

C ars, a ipla nes. . => c omplic a ted, not c omplex

Complex (no unique definition):

- many interacting units
- no c entra lized a uthority, self-o rga nized
- complic ated at all sc a les
- evolving structures
-emerging properties (hea vy-tails, hiera rchies...)

Exa mples: Intemet, WWW, Soc ial nets, etc ...

## Exa mple: Intemet growth



## Ma in features of c omplex networks

- Many interacting units
- Self-orga nization
-Small-world
- Scale-free heterogeneity
-Dyna mic al evolution


## Standard graph theory

Random graphs

- Sta tic
-Ad-hoc topology

Example: Intemet topology generators Modeling of the Intemet structure with ad-hoc algorithms tailored on the properties we consider more relevant

## Sta tistic al physic s a p proa ch

Mic roscopic processes of the many component units


Macroscopic statistical and dynamical properties of the system

Cooperative phenomena Complex topology

Natural outc ome of the dynamical evolution

## Robustness

## Complex systems ma inta in their basic functions

 even under emors and failures (cell $\rightarrow$ mutations; Intemet $\rightarrow$ router breakdowns)S: fraction of giant component


Fraction of removed vertices, f

vertex fa ilure


## Case of Scale-free Networks

## S <br> Random failure $\mathrm{f}_{\mathrm{c}}=1 \quad(2<\gamma \leq 3)$

Attack $=$ progressive failure of the most Connected vertices $\mathrm{f}_{\mathrm{C}}<1$
$\mathrm{f}_{\mathrm{c}} \quad 1$
Intemetmaps

R. Albert, H. J eong, A.L. Bara basi, Nature 406378 (2000)

## Failures vs, attacks

## Failures

Topological enrortolerance




## Fa ilures = perc olation

ffraction of vertic es removed because of failure
p=probability of a vertex to be present in a percolation problem

Question: existence or not of a giant/percolating cluster, i.e. of a connected cluster of nodes of size O(N)

## Betweenness

$\Rightarrow$ mea sures the "centrality" of a vertex i:
for each pair of vertices $(1, m)$ in the graph, there are $\sigma^{\mathrm{Im}}$ shortest paths between I and $m$ $\sigma_{i}^{l m}$ shortest paths going through i
$b_{i}$ is the sum of $\sigma_{i}^{\text {Im }} / \sigma^{\text {lm }}$ over all pairs $(\mathrm{l}, \mathrm{m})$

$b_{j}$ is la rge
$b_{j}$ is small

## Atta c ks: other strategies

* Most connected vertices
* Vertic es with largest betweenness
* Removal of edges linked to vertic es with large $k$
* Removal of edges with largest betweenness
* Cascades
P. Holme et al (2002); A. Motter et al. (2002);
D. Watts, PNAS (2002); Dall'Asta et al. (2006)...


## Classical random graphs

Solomonoff \& Rapoport (1951), Erdös-Rényi (1959)
$\mathcal{G}_{n, p}=\begin{aligned} & \mathrm{n} \text { vertices, edges with probability } \\ & \mathrm{p}: \text { static random graphs }\end{aligned}$


Average number of edges: $\varangle \mathbb{=}=\mathrm{pn}(\mathrm{n}-1) / 2$
Average degree: $<k>=p(n-1)$
$p=c / n$ to have
finite average degree
Related formulation
$\mathcal{G}_{n, m}=\mathrm{n}$ vertices, $m$ edges: each configuration is

## Classic al random graphs

Probability to have a vertex of degree $k$

- connected to $k$ vertices,
- not connected to the othern-k-1

$$
P(k)=\binom{n-1}{k} p^{k}(1-p)^{n-k-1}
$$

Large $N$, fixed $p N=<k>$ : Poisson distribution

$$
P(k)=e^{-\langle k\rangle} \frac{\langle k\rangle^{k}}{k!}
$$

Exponential decay at large $k$


## Classical random graphs

## Properties

I. Poisson degree distribution (on la rge graphs)
II. Small clustering coefficient: <C>=p=<k>/n (goesto zero in the limit of infinite graph size for sparse graphs)
III. Short distances: dia meter I ~log (n)/log( $\langle k\rangle$ ) (number of neighbors at distance $d$ : $<k>d$

## Classical random graphs

$<k><1$ : many small subgraphs
<k>>1: giant c omponent + small subgraphs


## Classical random graphs

Self-c consistent equation for relative size of giant component

$$
\begin{gathered}
u=\sum_{k=0}^{\infty} P(k) u^{k}=e^{-<k>} \sum_{k=0}^{\infty} \frac{(<k>u)^{k}}{k!}=e^{<k>(u-1)} \\
S=1-u \rightarrow S=1-e^{-<k>S}
\end{gathered}
$$

Avera ge cluster size

$$
\langle s\rangle=\frac{1}{1-<k>+<k>S}
$$

## Classical random graphs



Near the transition (<k>~1)

$$
\begin{gathered}
S \sim(<k>-1)^{\beta} \\
<s>\sim|<k>-1|^{-\gamma} \\
\beta=1, \gamma=1
\end{gathered}
$$

## Generalized random graphs

The configuration model (Molloy \& Reed, 1995, 1998)

Basic idea: building random graphs with a rbitrary degree distributions

How it works
I. Choose a degree sequence compatible with some given distribution
II. Assign to each vertex his degree, ta ken from the sequence, in that each vertex has as many outgoing stubs a sits degree
III. J oin the stubs in random pairs, until all stubs a re joined

## Generalized random graphs

Simple model: easy to handle a nalytically

Important point: the probability that the degree of vertex reached by following a randomly chosen edge is k is not given by $P(k)$ !

Reason: vertices with large degree have more edges and can be reached more easily than vertices with low degree

Conclusion: distribution of degree of vertic es at the end of randomly chosen edge is proportional to $\mathrm{KP}(\mathrm{k})$

## Generalized random graphs

Excess degree: number of edges leaving a vertex reached from a randomly selected edge, other than the edge followed

Distribution of excess degree: $Q(k)$

$$
Q(k)=\frac{(k+1) P(k+1)}{\sum_{k} k P(k)}=\frac{(k+1) P(k+1)}{<k>}
$$

## Generalized random graphs

Chance of finding loops within small (i.e. non-giant) components goesas $O\left(n^{-1}\right)$

Consequence: graphs of the configuration model are essentially loopless, tree-like, unlike real-world networks

Generating functions

$$
\begin{aligned}
& G_{0}(x)=\sum_{k=0}^{\infty} P(k) x^{k} \\
& G_{1}(x)=\sum_{k=0}^{\infty} Q(k) x^{k}
\end{aligned}
$$

## Generalized random graphs

$$
G_{1}(x)=G_{0}^{\prime}(x) /<k>
$$

$$
<k>=G_{0}^{\prime}(1)
$$

$$
<k^{2}>-<k>=G_{0}^{\prime}(1) G_{1}^{\prime}(1)
$$

$$
<s>=1+\frac{<k>^{2}}{2<k>-<k^{2}>}
$$

Condition for existence of giant component:

$$
\sum_{k} k(k-2) P(k)=0 \quad \leftrightarrow \quad G_{1}^{\prime}(1)=1
$$

## Generalized random graphs

 $u=$ probability that a randomly chosen edge is not in the giant component$$
\begin{gathered}
u=\sum_{k=0}^{\infty} Q(k) u^{k}=G_{1}(u) \\
1-S=\sum_{k=0}^{\infty} P(k) u^{k}=G_{0}(u)
\end{gathered}
$$

Self-c onsistent equations for the rela tive size S of the giant component

## Generalized random graphs

Example: power-law degree distribution

$$
P(k)= \begin{cases}0 & \text { for } k=0 \\ k^{-\alpha} / \zeta(\alpha) & \text { for } k \geq 1\end{cases}
$$

$G_{0}(x)=\frac{L i_{\alpha}(x)}{\zeta(\alpha)}, \quad G_{1}(x)=\frac{L i_{\alpha-1}(x)}{x \zeta(\alpha-1)}$

$$
\operatorname{Li}_{s}(z)=\sum_{k=1}^{\infty} \frac{z^{k}}{k^{s}}
$$

$$
G_{1}^{\prime}(1)=1 \rightarrow \zeta(\alpha-2)=2 \zeta(\alpha-1)
$$

Critical exponent value: $\alpha_{c}=3.4788$
For $\alpha<\alpha_{c}$ there is always a giant connected For aponentiere is no giant connected component!

## Generalized random graphs

$$
S=1-G_{0}(u), \quad u=G_{1}(u)
$$

$$
u=\frac{\operatorname{Li}_{\alpha-1}(u)}{u \zeta(\alpha-1)}
$$

For $\alpha<2, \mathrm{u}=0$ and $\mathrm{S}=1 \Rightarrow$ All vertices are in the giant component!
W. Aiello, F. Chung, L. Lu, Proc. 32th ACM Symposium on Theory of Computing 171-180 (2000)

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vI. Dynamic processes in networks

